Network Analysis:

The Hidden Structures behind the Webs We Weave 17-338 / 17-668

Power and Centrality in Social Networks

Tuesday, September 23, 2025

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"Central" actors are powerful

Relational view of the individual

So far, we covered dyads and triads as basic building blocks of networks

The relational view has been the overarching theme

- relations between two people
- relations among three people

The same relational view to study the individual

This approach sharply contrasts with widely used approaches

- Surveys ask individuals about their attributes, beliefs, and actions
- Majority of statistical models assume that individuals are like disconnected atoms
 - Independent, identically distributed variables assumption (IID).

Which node is the most "prominent"?

In contrast, network analysis views an individual in relation to their network neighbors

- Opinions are influenced by network neighbors
- Individual's **status** and importance comes from their position in the network
- Individuals can leverage their position to influence opinions, selectively spread or block information, and control opportunities
- **Power** that an individual can wield is fundamentally relational

A network position of power and prominence is a "central" position

- Power and prominence are multifaceted
- Therefore, "central" can reflect different aspects of power
 - "Centrality" can be defined and measured in different ways

So many ways to define "centrality"

group_closeness_centrality (G, S[, weight])

group degree centrality (G, S)

group_in_degree_centrality (G, S)

group_out_degree_centrality (G, S)

prominent_group (G, k[, weight, C, ...])

(Shortest Path) Betweenness

betweenness centrality (G[, k, normalized, ...])

betweenness_centrality_subset (G, sources, ...)

edge_betweenness_centrality (G[, k, ...])

edge_betweenness_centrality_subset (G, ...

[, ...])

Compute the shortest-path betweenness

centrality for nodes.

subset of nodes.

Compute betweenness centrality for edges.

Compute betweenness centrality for edges

for a subset of nodes.

Compute betweenness centrality for a

Centrality		Current Flow Betweenness		Load	Trophic
Degree degree_centrality (G) Compute the degree centrality for nodes.		${\tt current_flow_betweenness_centrality} \; ({\tt Gf}, \ldots))$	Compute current-flow betweenness centrality for nodes.	load_centrality (G[, v, cutoff, normalized,]) Compute load centrality for nodes. edge_load_centrality (G[, cutoff]) Compute edge load.	<pre>trophic_levels (G[, weight]) trophic_differences (G[, weight])</pre>
degree_centrality (G) Compute the degree centrality (G) Compute the in-degree centrality (G) Compute the out-degree centrality (G) Compute the degree centrality	entrality for nodes.	<pre>edge_current_flow_betweenness_centrality (G)</pre>	Compute current-flow betweenness centrality for edges.	Subgraph subgraph_centrality (G) Returns subgraph centrality for each node in G.	trophic_incoherence_parameter (G[, weig
Eigenvector	·	approximate_current_flow_betweenness_centrality (G)	Compute the approximate current-flow betweenness centrality for nodes.	subgraph_centrality_exp (G) Returns the subgraph centrality for each node of G. estrada_index (G) Returns the Estrada index of a the graph G.	VoteRank voterank (G[, number_of_nodes]) Selection
eigenvector_centrality (G[, max_iter, tol,]) Compute the eigenvector centrality numbur (G[, weight,])	eigenvector centrality for the graph 6. eigenvector centrality for the graph G.	current_flow_betweenness_centrality_subset (G,)	Compute current-flow betweenness centrality for subsets of nodes.	Harmonic Centrality [harmonic_centrality (G[, nbunch, distance,]) Compute harmonic centrality for nodes.	
katz_centrality (G[, alpha, beta, max_iter,])	atz centrality for the nodes of the graph G. atz centrality for the graph G.	edge_current_flow_betweenness_centrality_subset (G,)	Compute current-flow betweenness centrality for edges using subsets of nodes.	Dispersion	Laplacian Laplacian_centrality (G[, normalized,
Closeness closeness_centrality (G[, u, distance,]) incremental_closeness_centrality (G, edge[,]) Incremental_closeness_centrality (G, edge[,])	oseness centrality for nodes.	Communicability Betweenness communicability betweenness centrality (G) Returns subg	graph communicability for all pairs	Reaching [local_reaching_centrality (G, v[, paths,]) Returns the local reaching centrality of a node in	
Current Flow Closeness			of nodes in G.	a directed graph. Returns the global reaching centrality (G[, weight,]) directed graph.	
information centrality (Gf, weight dtype]) Compute current-fl	nt-flow closeness centrality for nodes. nt-flow closeness centrality for	Group Centrality [group_betweenness_centrality (G, C[,])] Compute the g	group betweenness centrality for a	Percolation percolation centrality (Gf. attribute)) Compute the percolation centrality for nodes.	

Compute the group closeness centrality for a

group of nodes.

Compute the group degree centrality for a group

of nodes.

Compute the group in-degree centrality for a

group of nodes.

Compute the group out-degree centrality for a

group of nodes.

Find the prominent group of size k in graph G.

second_order_centrality (G) Compute the second order centrality for nodes of G.

Second Order Centrality

So many ways to define "centrality"

We will cover the four most widely used centrality measures in network analysis

- Degree centrality
- Closeness centrality
- Betweenness centrality
- Eigenvector centrality

Getting Part the Gatekeeper:



Insight: Central nodes have many connections

- Access to information
 - Gate keepers



Insight: Central nodes have many connections

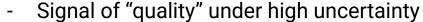
- Access to information
 - Gate keepers

- Concentration of attention (attention is scarce)
 - Opinion leaders

Insight: Central nodes have many connections

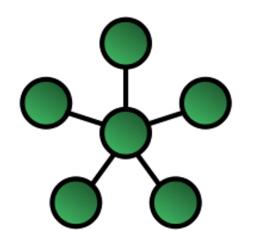
- Access to information
 - Gate keepers

- Concentration of attention (attention is scarce)
 - Opinion leaders



- Influencers
- 'There must be some reason why he has so many connections.'





Insight: Central nodes have many connections

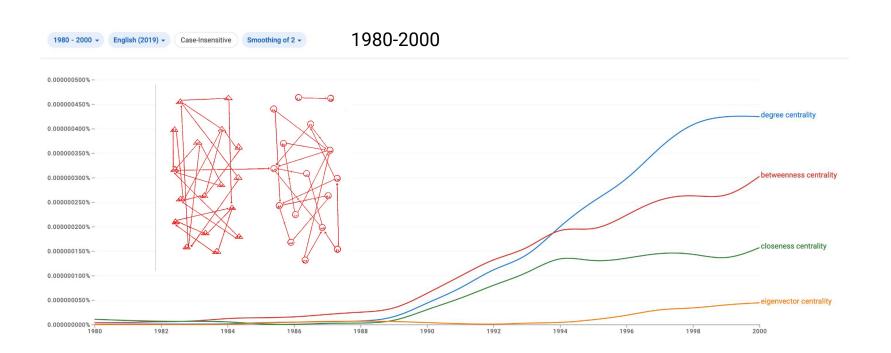
The number of ties of a node, **normalized by network** size (g-1).

$$C_D(i) = \frac{k_i}{g-1}$$

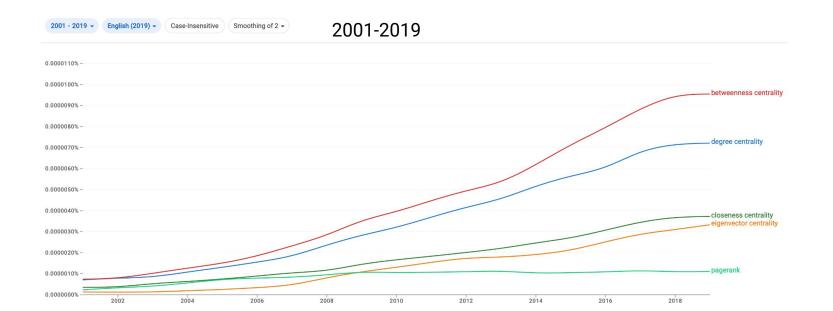
Normalization allows for comparison across different networks

One of the most basic and fundamental quantities in network science

Degree centrality was, by far, the most popular measure up to the early 2000s



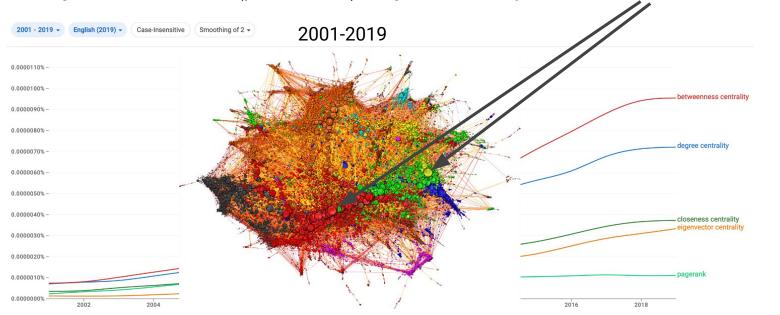
However, its prominence arguably diminished in the past two decades **Why?**



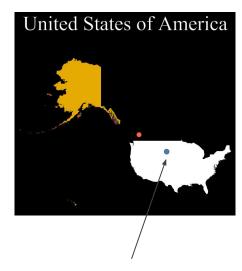
However, its prominence arguably diminished in the past two decades

Why?

In large-scale networks (post-2000s), degree centrality becomes a "local" measure



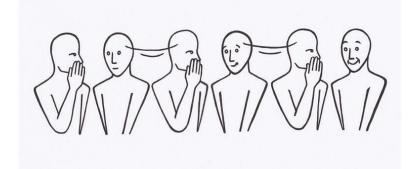
Closeness Centrality



Centroid of the U.S.

Insight: The position that minimizes the distance to all other positions is the most central

Closeness Centrality



Telephone game

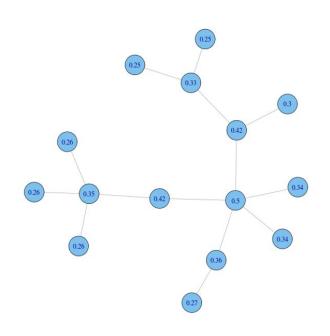
Insight: The position that minimizes the distance to all other positions is the most central

Application:

- Systems where traversing the network is costly
- Information attrition rate is high

- → The node closest to all other nodes can obtain more accurate messages
- → Faster access
- → A source of power

Closeness Centrality



Central nodes have short paths to other nodes

Calculated as the reciprocal of the sum of pairwise path distances

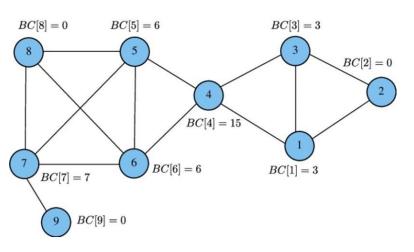
$$C_c(i) = [\sum_{j=1}^g d(i,j)]^{-1}$$

Normalized by size of network

$$C_c(i) = (g-1)[\sum_{j=1}^g d(i,j)]^{-1}$$

A global measure that uses information from the entire network

Number of shortest paths node is on

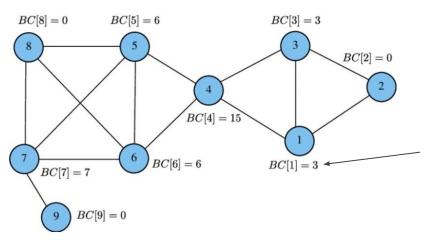


Node is central if it sits on many shortest paths

Nodes positioned at Information bottlenecks are central

Those nodes have more control over the distribution of information and other resources

Number of shortest paths node is on

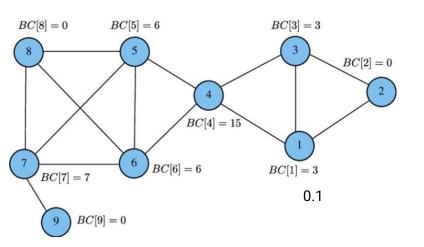


Node *i*'s betweenness is the sum of the probabilities that *i* is on the shortest paths of all node pairs in the network

$$C_B(i) = \sum_{j < k} PL(i, j, k) / PL(j, k)$$

 $PL(i,j,k) \rightarrow$ Number of shortest paths involving i

 $PL(j,k) \rightarrow$ Number of shortest paths between j and k

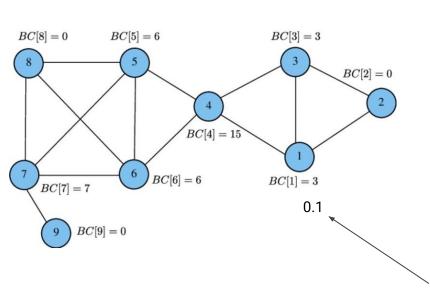


$$C_B(i) = \sum_{j < k} PL(i, j, k) / PL(j, k)$$

Maximum probability PL(i,j,k) / PL(j,k) = 1 is when i sits on every shortest path between j and k.

Then, the theoretical maximum C(i) is when i sits on every shortest path for every pair of nodes (excluding i)

Q: What is the number of node pairs (dyads) excluding *i* in a graph with *g* nodes?



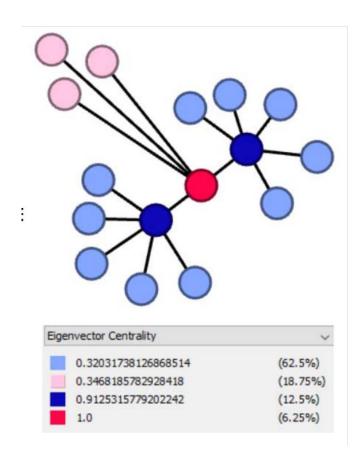
$$C_B(i) = \sum_{j < k} PL(i, j, k) / PL(j, k)$$

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Then, the theoretical maximum C(i) is when i sits on every shortest path for every pair of nodes (excluding i)

 \rightarrow Normalize by the total number of node pairs excluding $i \rightarrow (g-1)(g-2)/2$

$$C_B'(i) = \frac{C_B(i)}{\left[\frac{(g-1)(g-2)}{2}\right]}$$

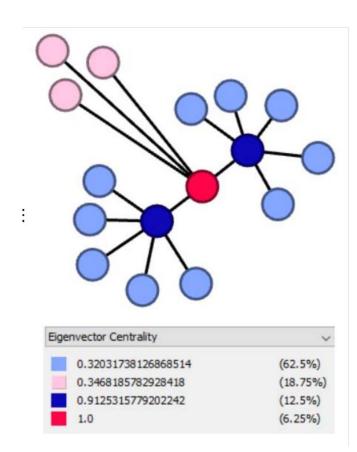


Who is more central?

→ Someone who knows five (well connected) celebrities vs. someone who knows five ordinary people

Eigenvector centrality quantifies this insight

→One's Eigenvector centrality is determined by the neighbors' eigenvector centrality, which in turn are determined by their neighbors' eigenvector centrality ...



Who is more central?

 \rightarrow Someone who knows five (well connected) celebrities vs. someone who knows five ordinary people

Eigenvector centrality quantifies this intuition

→One's Eigenvector centrality is determined by the neighbors' eigenvector centrality, which in turn are determined by their neighbors' eigenvector centrality ...

Red node has two (purple) friends who are themselves highly connected

→ Highest Eigenvector centrality

Pink nodes have same degree centrality as the sky blue nodes, but they are connected to the red node (highest Eigenvector centrality)

→ Eigenvector centrality: pink > sky blue

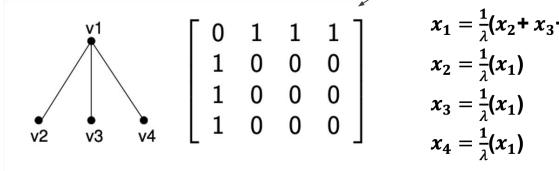
Given an adjacency matrix **A**, eigenvector centrality of node *i* is:

i's weighted degree where each of i's edge is weighted by the degree of neighbor, j

$$\lambda C_E(i) = \sum_j A_{ij} C_E(j)$$

Given adjacency matrix ${\pmb A}$, compute the eigenvector ${\pmb x}$ corresponding to the principal eigenvalue ${\pmb \lambda}^*$ of ${\pmb A}$

Each element in **x** is the eigenvector centrality value of the corresponding node



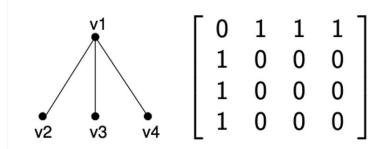
For adjacency matrix A: Eigenvector centrality of the nodes are

$$x_1 = \frac{1}{\lambda}(x_2 + x_3 + x_4)$$

$$x_2 = \frac{1}{\lambda}(x_1)$$

$$x_3 = \frac{1}{\lambda}(x_1)$$

$$x_4 = \frac{1}{\lambda}(x_1)$$



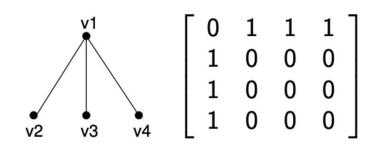
$$\lambda x_1 = (x_2 + x_3 + x_4)$$

$$\lambda x_2 = x_1$$

$$\lambda x_3 = x_1$$

$$\lambda x_4 = x_1$$

$$\lambda \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix}$$



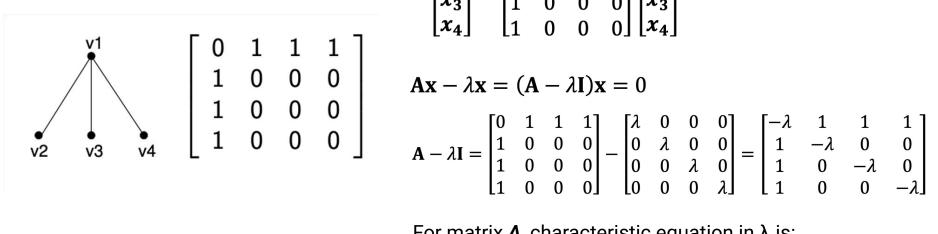
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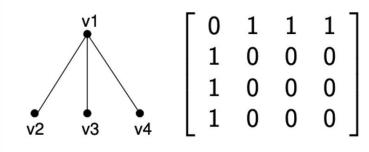


$$\lambda \begin{vmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{vmatrix} = \begin{vmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{vmatrix} \begin{vmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{vmatrix} \longrightarrow \lambda x = Ax$$

$$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 0 \end{bmatrix} \begin{bmatrix} \lambda & 0 & 0 & 0 \\ 0 & \lambda & 0 & 0 \end{bmatrix} \begin{bmatrix} -\lambda & 1 & 1 & 1 \\ 1 & -\lambda & 0 & 0 \end{bmatrix}$$

For matrix **A**, characteristic equation in λ is:

$$\begin{bmatrix} -\lambda & 1 & 1 & 1 \\ 1 & -\lambda & 0 & 0 \\ 1 & 0 & -\lambda & 0 \\ 1 & 0 & 0 & -\lambda \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$



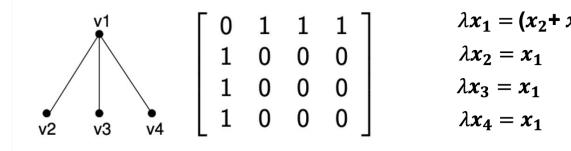
$$\lambda \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} \longrightarrow \lambda x = Ax$$

Solution to the characteristic equation:

 \rightarrow determinant of $\mathbf{A} - \lambda \mathbf{I}$ is $\mathbf{0}$

$$\lambda^4 - 3\lambda^2 = \lambda^2(\lambda^2 - 3) = 0$$

Eigenvalues are $-\sqrt{3}$, 0, 0, $\sqrt{3}$ Principal eigenvalue $\lambda^* = \sqrt{3}$



$$\lambda x_1 = (x_2 + x_3 + x_4)$$
 $\lambda x_2 = x_1$
 $\lambda x_3 = x_1$
 $\lambda x_4 = x_1$
 $\sqrt{3}x_1 = x_2 + x_3 + x_4$
 $\sqrt{3}x_2 = x_1$
 $\sqrt{3}x_3 = x_1$
 $\sqrt{3}x_4 = x_1$

$$\mathbf{x} = \begin{bmatrix} \sqrt{3} \\ 1 \\ 1 \\ 1 \end{bmatrix} \quad \text{v1 has highest eigenvector}$$

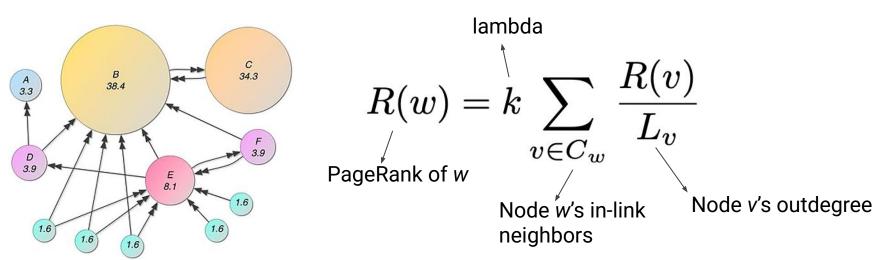
$$\text{centrality, } \sqrt{3}$$

$$\text{while } \mathbf{v2} = \mathbf{v3} = \mathbf{v4} = 1$$

$$\lambda x = Ax \longrightarrow \sqrt{3} \begin{bmatrix} \sqrt{3} \\ 1 \\ 1 \\ 1 \end{bmatrix} = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix} \begin{bmatrix} \sqrt{3} \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

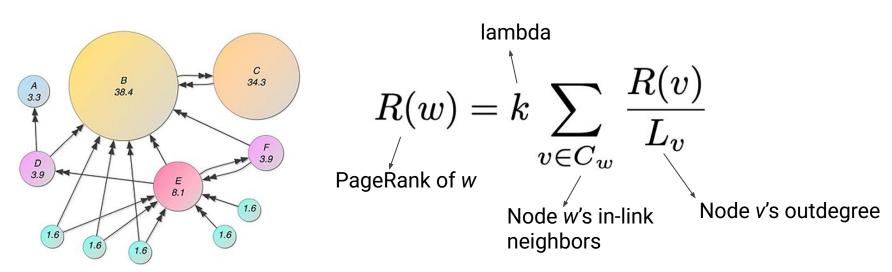
Google's original PageRank algorithm ranks web pages using a variant of eigenvector centrality

- A node's PageRank is a function of the PageRank of its in-neighbors
- Normalize the eigenvector centrality of the in-neighbor nodes with their outdegree



Google's original PageRank algorithm ranks web pages using a variant of eigenvector centrality

- A node's PageRank is a function of the PageRank of its in-link neighbors
- Normalize the eigenvector centrality of the in-link neighbors with their outdegrees
 - Why?



Similarities among Centrality Measures

Average correlations between centrality measures (N=58).

Pearson correlation

		1	2	3	4	5	6	7	8	9	10	11	
1	Indegree												
2	Outdegree	0.3											
3	Degree	0.78	0.71										
4	Between	0.62	0.54	0.7	,								
5	S-Between	0.69	0.5	0.85	0.67								
6	Closeness-In	0.55	0.16	0.45	0.37	0.31							
7	Closeness-Out	0.18	0.81	0.56	0.39	0.38	0.02						
8	S-Closeness	0.4	0.64	0.66	0.37	0.44	0.42	0.65					
9	Integration	0.7	0.26	0.58	0.5	0.41	0.9	0.15	0.51				
10	Radiality	0.21	0.86	0.61	0.44	0.41	0.06	0.98	0.67	0.19			
11	S-Int/Rad	0.45	0.7	0.73	0.43	0.5	0.44	0.69	0.99	0.54	0.72		
12	Eigenvector	0.71	0.69	0.92	0.64	0.72	0.44	0.55	0.63	0.57	0.59	0.71	
Average		0.51	0.56	0.69	0.52	0.53	0.37	0.49	0.58	0.48	0.52	0.63	0.65
Standard	Deviation	0.21	0.23	0.14	0.16	0.14	0.27	0.22	0.25	0.28	0.17	0.12	0.12

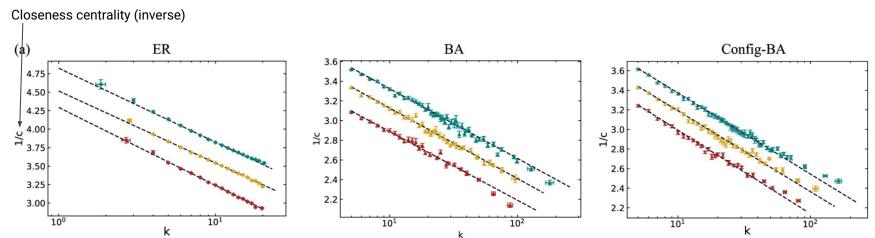
** **S**-Between: Betweenness centrality on **s**ymmetrized network

Valente et al. 2008

Similarities among Centrality Measures

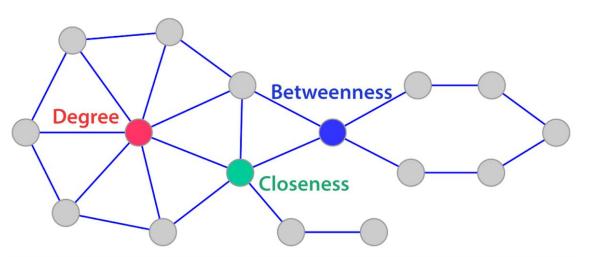
Sometimes, the relationship is non-linear \rightarrow lower linear correlation (Pearson)

- Example: Closeness centrality and the log of degree centrality are highly correlated (redundant)
- If such a strong relationship, why do we care about computationally expensive closeness centrality?



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Similar, but also different



Degree v	ClosenessCentrality	BetweennessCentrality
7	0.45454545	0.29047619
5	0.51724138	0.42380952
4	0.48387097	0.4952381

Different centrality measures capture different intuitions

In a graph, the node with highest degree is not necessarily the node with the highest betweenness or closeness

Summary

Centrality as a general term for measuring a node's position of "prominence" in the network

Degree: sheer connections

Closeness: shortest path distance

Betweenness: node's presence in shortest paths

Eigenvector: degree weighted by the degree of the neighbors

These measures are highly correlated, but conceptually distinct